

Don't Abandon the Primary Key: A High-Synchronization and Robust Virtual Primary Key Scheme for Watermarking Relational Databases

Ke Yang^{1,2,3} Shuguang Yuan^{1,2,3} Jing Yu^{1,2,3} Yuyang Wang^{1,2,3} Tengfei Yang⁴ Chi Chen^{1,2,3}

¹ Institute of Information Engineering, Chinese Academy of Sciences

² School of Cyber Security, University of Chinese Academy of Sciences

³ Key Laboratory of Cyberspace Security Defense

⁴ National Computer Network Emergency Response Technical Team/Coordination Center of China



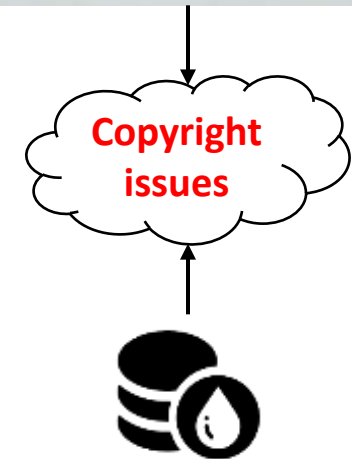
Background

- Relational databases are still the data storage and sharing solution in most enterprises and applications.
- The demand for databases publishing and sharing has increased significantly.
- Copyright issues become a significant concern.
 - Unauthorized database distribution/modification
 - Database piracy



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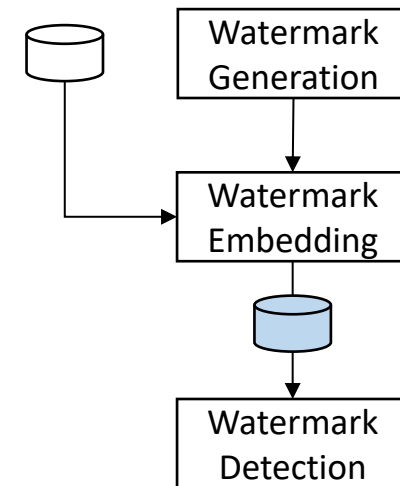
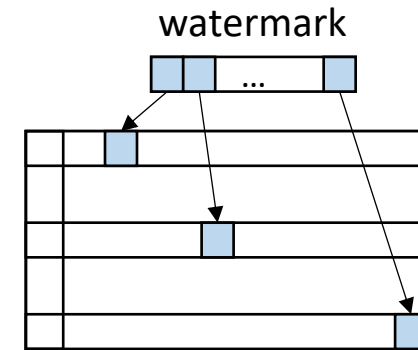
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 - Unauthorized database distribution/modification
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- **Database Watermarking techniques** can effectively protect copyright.



Background

What is database watermarking and how it work?

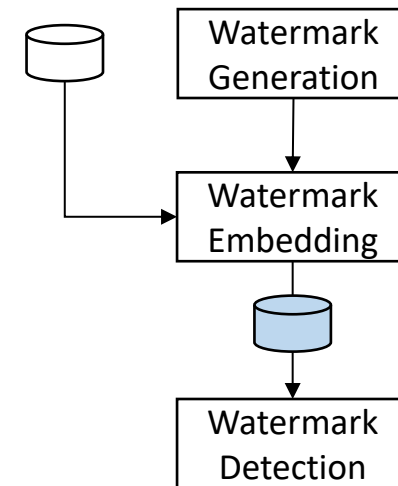
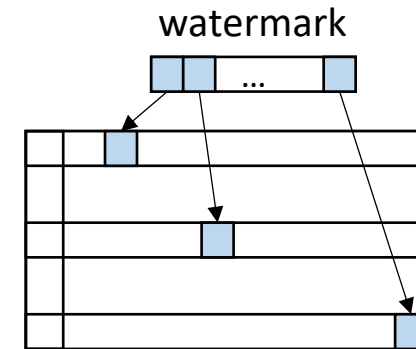
- Database watermarking secretly embeds hidden markers or watermark information within a database to protect copyright.
- Workflow of database watermarking:
 - Watermark Generation
 - Watermark Embedding
 - Watermark Detection



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- Workflow of database watermarking:
 - Watermark Generation
 - Watermark Embedding
 - Watermark Detection
- Main goals: Prevent copyright issues, and provide evidence of databases ownership by **detecting watermarks** in case of disputes.



Background

Watermark synchronization

The process of aligning the detected watermark with the embedded watermark.

PK scheme

- Take advantage of the uniqueness of primary key
- Achieve high watermark synchronization

Vulnerability

- Primary key attacks: Cannot detect watermark when primary key is erased or changed.

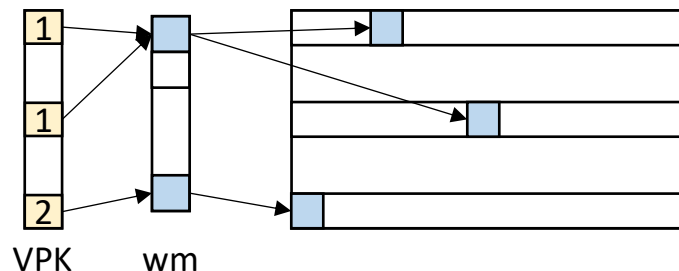
Background

VPK scheme

- Generate virtual primary key for watermarking instead of primary key
- Address primary key attacks

Vulnerability

- Synchronization problems: The redundant values of virtual primary keys affect the embedding quality and watermark synchronization.
- Fail to utilize primary key: Even primary key remains unchanged, which decrease detection ability in many cases.



Background

Robust against attacks

- **Subset attacks** (distortion of tuples)
- **Attribute attacks** (distortion of attribute)
- **Additive attacks** (re-watermarking)

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New risk we found

- **Attribute name attacks:** Attacks on attribute names (distortion of attribute names)
- Attribute name attacks can compromise watermark detection by **distorting the link** between detected and watermarked attributes.

Research Objectives & Contributions

Objectives

- Resist attribute name attacks
- Avoid primary key attacks & Mitigate synchronization problems
- Utilize primary key if it is intact
- Evaluation

Contributions

- A attribute classifier to maintain the correctness of attributes
- A high-synchronization and robust VPK scheme
- Watermarking algorithm integrates primary key
- The performance is verified by experiments

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<i>PK</i>	<i>A</i> ₁	<i>A</i> ₂	<i>A</i> ₃	<i>A</i> ₄	<i>A</i> ₅
001	43	193	2	710	433
002	15	358	3	549	487
003	21	670	8	967	404
004	39	454	2	828	469

(a) PK scheme

<i>PK</i>	<i>A</i> ₁	<i>A</i> ₂	<i>A</i> ₃	<i>A</i> ₄	<i>A</i> ₅
001	43	193	2	710	433
002	17	358	3	549	486
003	21	672	8	967	404
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(b) VPK scheme

<i>PK</i>	<i>A</i> ₁ <i>A</i> ₂ <i>A</i> ₃ <i>A</i> ₄ <i>A</i> ₅				
	Classifier				
	<i>L</i> ₁	<i>L</i> ₂	<i>L</i> ₃	<i>L</i> ₄	<i>L</i> ₅
001	43	192	2	710	432
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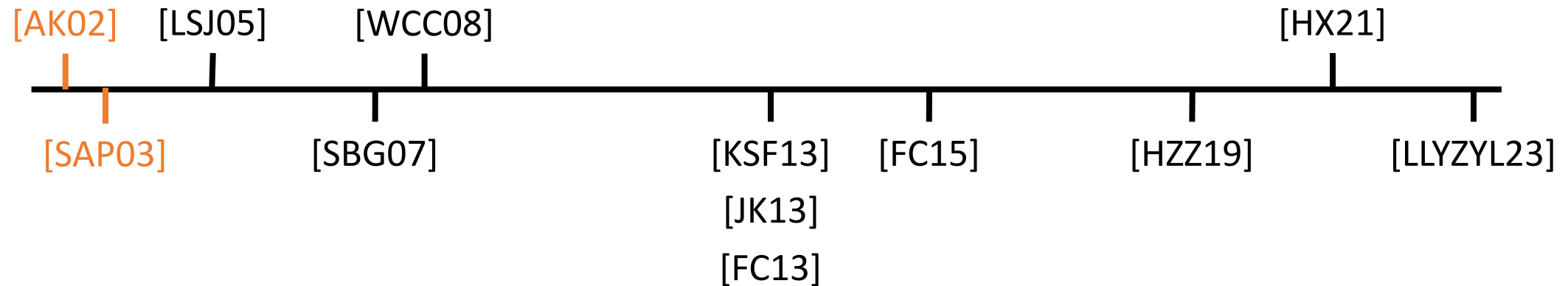
(c) Our scheme

Related Works

PK schemes

Bit-level Watermarking

Robust, meaningless watermark → Fingerprint, meaningful watermark, reversibility



Statistics-based → Distortion control: Optimization, DE, HS, semantic-preserved

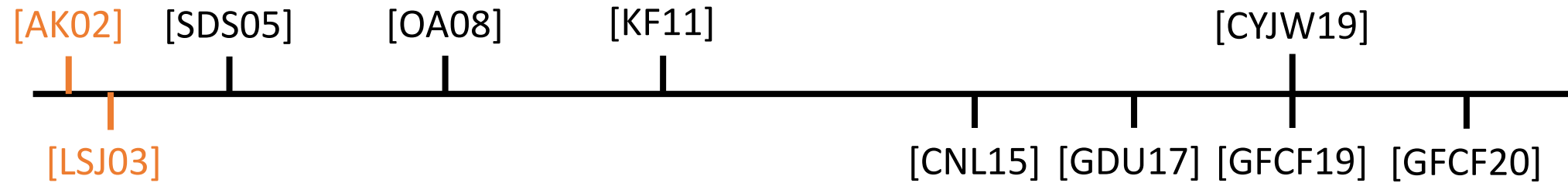
Statistics/semantic-based Watermarking

Related Works

VPK schemes

Non-primary key attributes as virtual primary key

Focus on statistics preserving and alternative attributes selecting



Focus on non-numerical relations and mitigate synchronization problems

Pseudo-random selected values in tuples to generate virtual primary key

Related Works

Attacks definition

- [KF18]
 - Insertion attacks, Deletion attacks, Alteration attacks
 - Multifaceted attacks
 - Additive (Re-watermarking) attacks
- [KSY20]
 - Substitution, addition, alteration
 - Vertical partition
 - Mix-Match
- [YCYYY22]
 - Subset attacks (Insertion attacks, Deletion attacks, Alteration attacks)
 - Attribute attacks

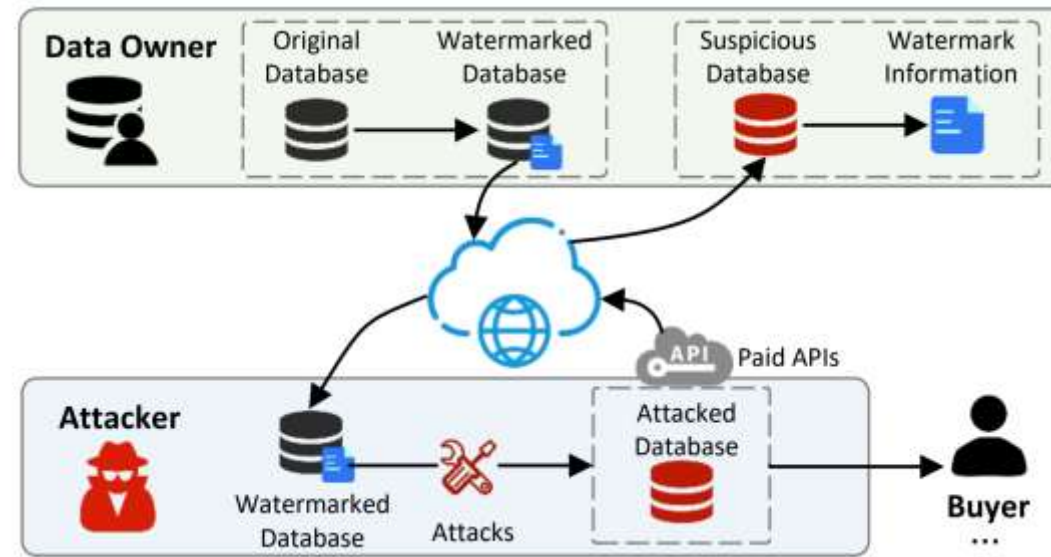
Threat Model

Defender

- Assumption: Watermark database + Detect watermark from suspicious database.

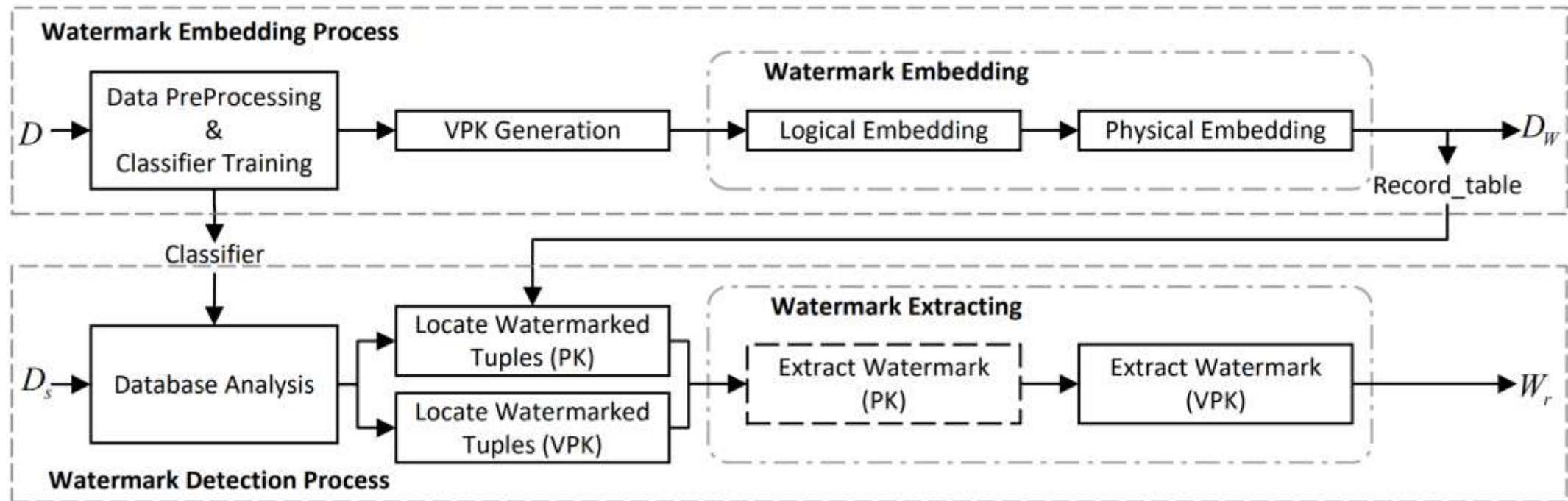
Attacker

- Goals: Bypass watermark detection + Maintain required data quality
- Capabilities: Access watermarking scheme + Alter database by attacks



Approach Overview

Workflow of watermarking scheme



Attribute Attacks Extension

Attribute column attacks

- Traditional attribute attacks: column deletion, insertion, or shuffling

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- Compromise attribute order
- Only [LSJ03] and [GFCF19] can resist

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Mix-match attribute attacks

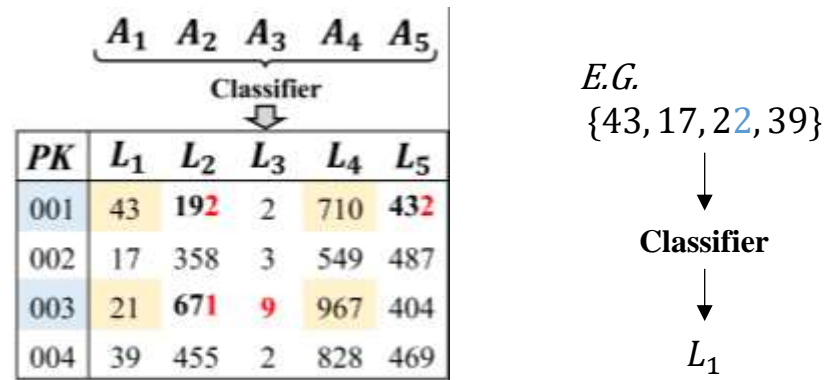
- Combinations of attribute column attacks and attribute name attacks
- Compromise attribute order & count

Data Preprocessing and Classifier Training

Q: How to resist Attribute name attacks and Mix-match attribute attacks?

Labeling attribute names and train a classifier to identify attribute labels.

Features: minimum, maximum, mean, variance, skewness, kurtosis, entropy of attributes.

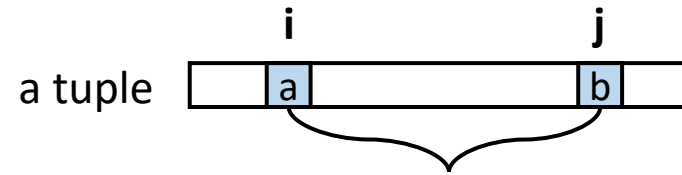


Accuracy: 96%, τ on prediction probabilities to trade off FP and FN.

Virtual Primary Key Generation

Candidate Attributes with Labels

- Selected to generate virtual primary key
- Decrease duplicate virtual primary key



vpk: Hash Combination $[H(a \cdot i), H(b \cdot j), H(a \cdot i \ b \cdot j)]$

vpk_dict: $a \cdot i \begin{cases} H(a \cdot i) \\ H(a \cdot i \ b \cdot j) \end{cases} \quad b \cdot j \begin{cases} H(b \cdot j) \\ H(a \cdot i \ b \cdot j) \end{cases}$

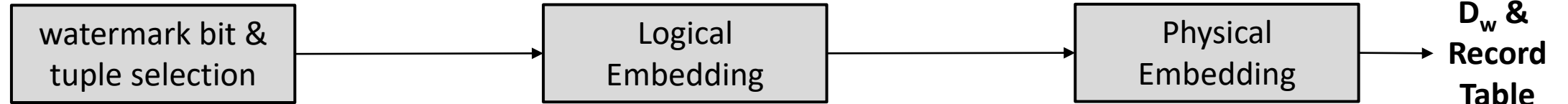
A Candidate Attribute value map to multiple virtual primary keys

- Increase watermark capacity & Mitigate synchronization problems
- Resist attribute column attacks

Watermark Embedding

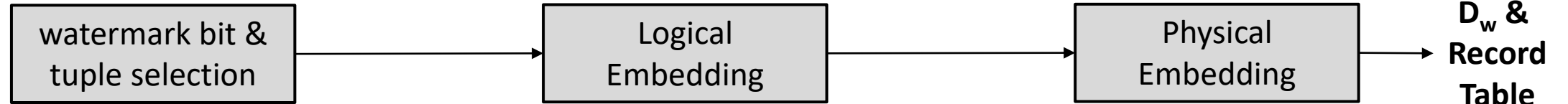


Watermark Embedding



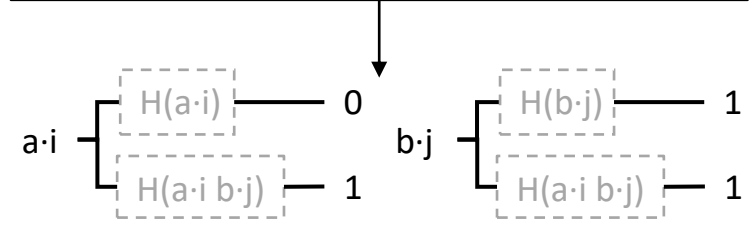
- PK select a watermark bit W_{pk}
 $H(K_s \cdot PK) \bmod \zeta$
- PK select a tuple to be marked
 $H(K_s \cdot PK) \bmod \gamma == 0$

Watermark Embedding

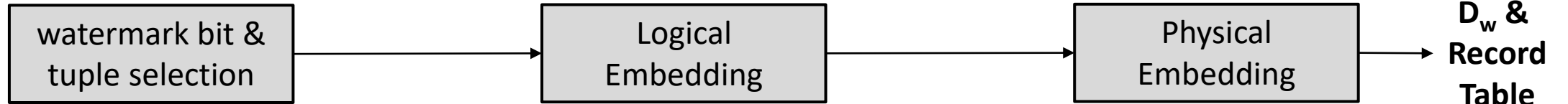


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- vpk select a watermark bit W_{vpk}
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- Compare $W_{pk} == W_{vpk}$, and record

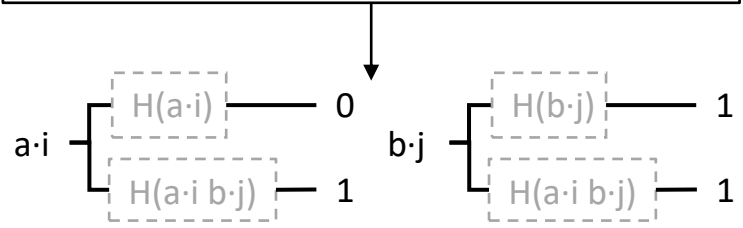


Watermark Embedding



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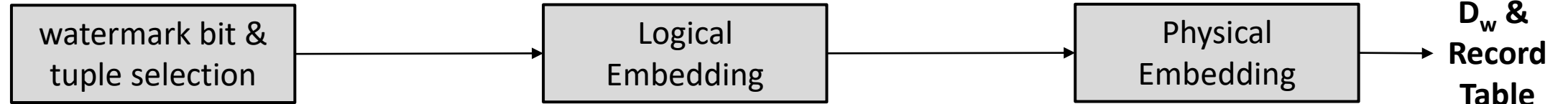
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- Select two attributes A_k, A_l
- set a bit to embed W_{pk}
 $H(K_s \cdot Len(A_k \ or \ l)) \bmod \xi$
- $ch = Ture, A_{k \ or \ l}$ changed

$A_k \cdot k = True, A_l \cdot l = False$

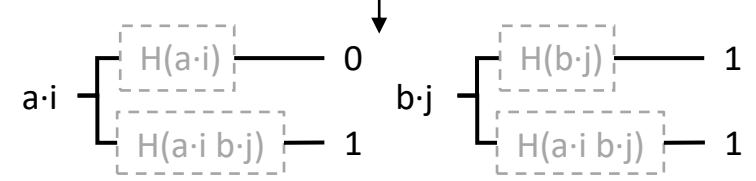
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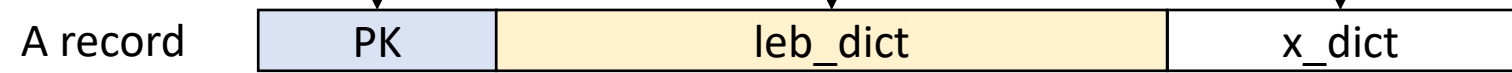
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Watermark Detection

Database Analysis

- Attack analysis, damage analysis
- Analysis of whether the label needs to be restored by attribute classifier

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Match record when Primary Key is Available

Matched PK	A_{w_i} of x_dict	$t.A_{w_i}$ of x_dict	$t.A_{w_i}$ of tuple t	Located water-marked tuple
9	A_{w_0}, A_{w_1}	7,3	7,3	True
12	A_{w_0}, A_{w_1}	24,5	22,7	False
37	A_{w_0}, A_{w_1}	13,2	13,3	True

$$wr = b$$

Extract watermark
by PK & VPK

W_r

$$\begin{cases} w_r = b, & \text{if } leb[a \cdot i] = 0 \\ w_r = \neg b, & \text{if } leb[a \cdot i] = 1 \end{cases}$$

Match record when Primary Key is Unavailable

$\Phi_{C'}$	Φ_C of leb_dict	A_{w_i} of x_dict	$t.A_{w_i}$ of x_dict	$t.A_{w_i}$ of tuple t	Located water-marked tuple
$234 \circ C_i$	$234 \circ C_i, 32 \circ C_j$	A_{w_0}, A_{w_1}	7,3	7,3	True
$256 \circ C_i$	$256 \circ C_i, 54 \circ C_j$	A_{w_0}, A_{w_1}	24,5	22,7	False
$197 \circ C_i$	$197 \circ C_i, 69 \circ C_j$	A_{w_0}, A_{w_1}	13,2	13,3	False

Extract watermark
by VPK

W_r

Experiments

Dataset: Forest Cover Type (581012 tuples)

Control: around 2900 watermarked tuples

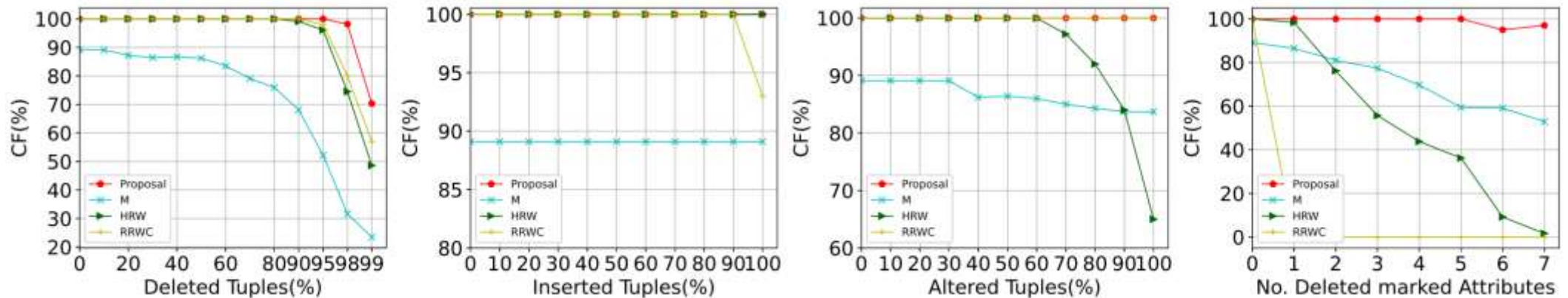
Robustness Verification:

VPK watermarking algorithm

Synchronization problems mitigation

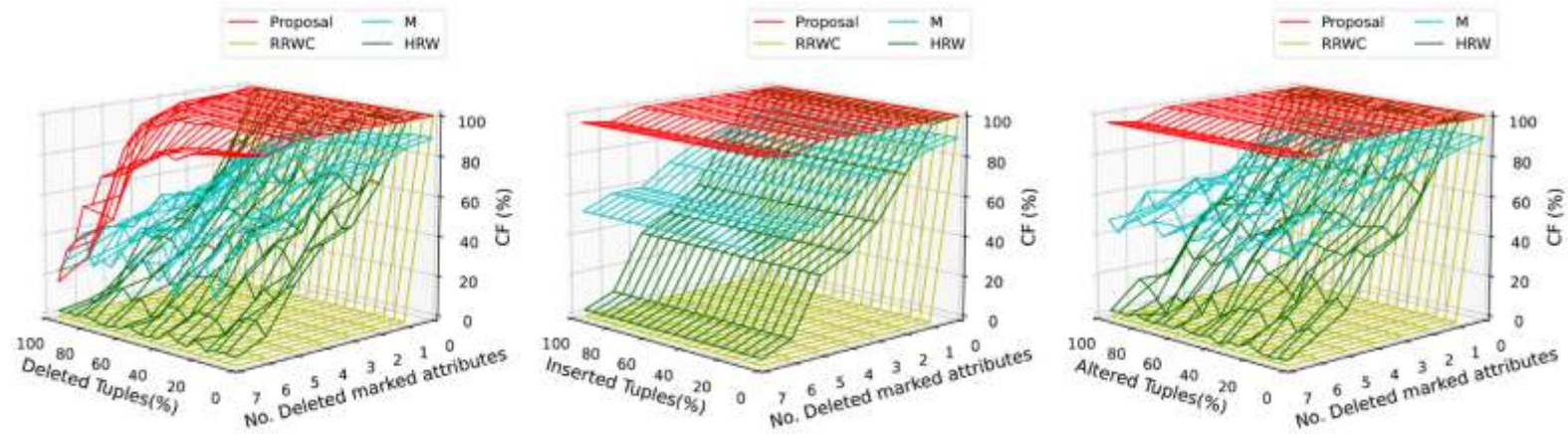
Algorithm	VPK	Unique VPK	ω_{max}	ω_{min}
Proposal	34761	3340	2331	689
M	3191	83	1326	0
HRW	2872	1181	161	33
RRWC	2906	2904	78	78

Normal single attack robustness verification

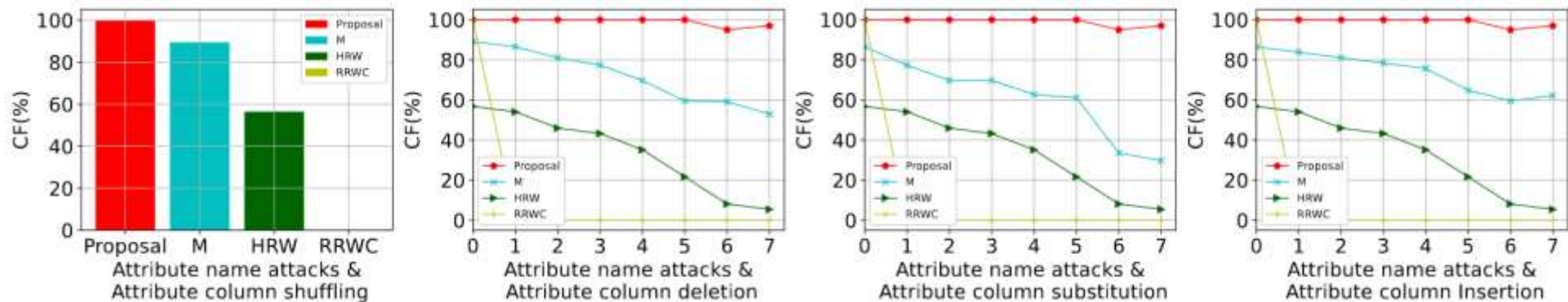


Experiments

Multifaceted attack robustness verification



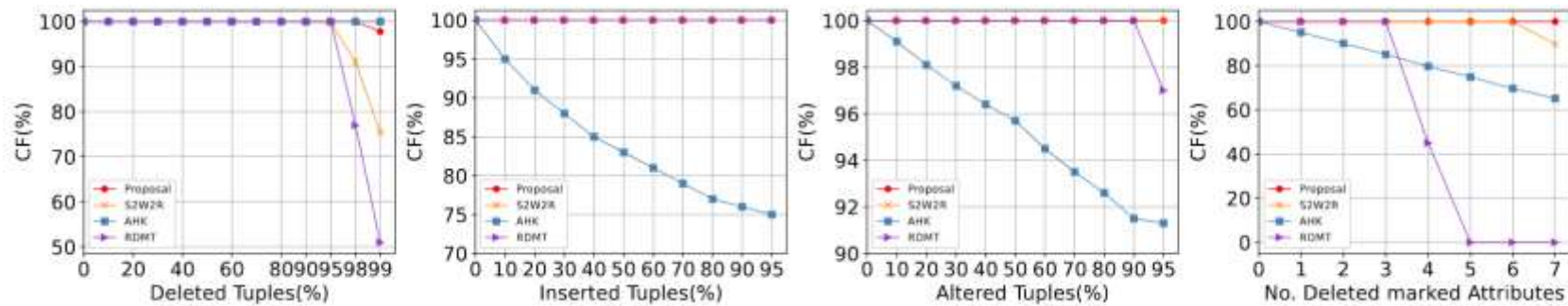
Mix-match attribute attack robustness verification



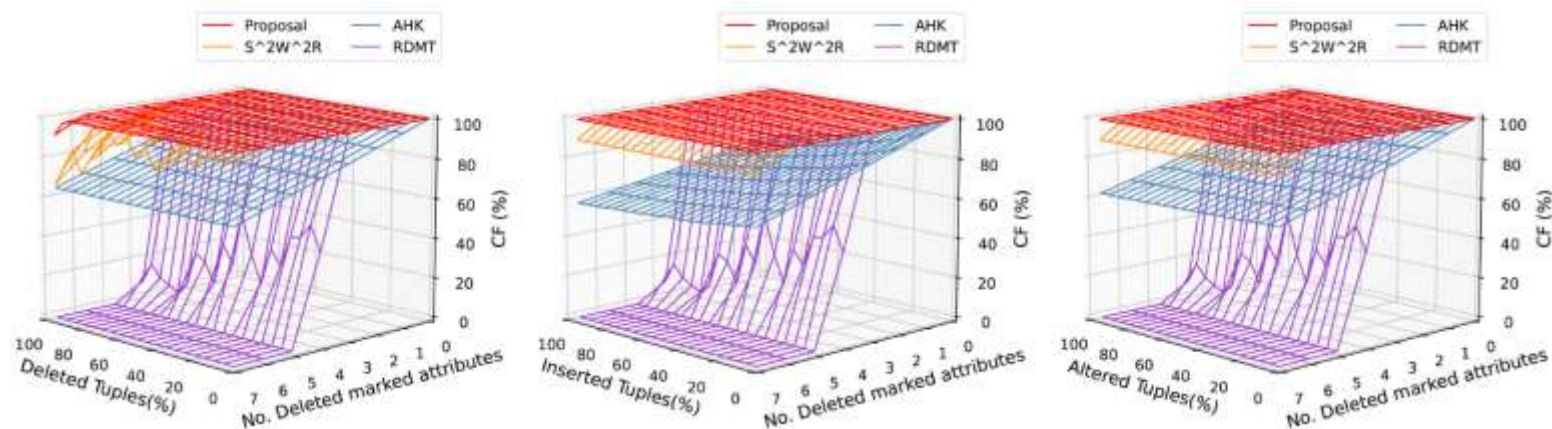
Experiments

Robustness Verification: PK watermarking algorithm

Normal single attack robustness verification



Multifaceted attack robustness verification



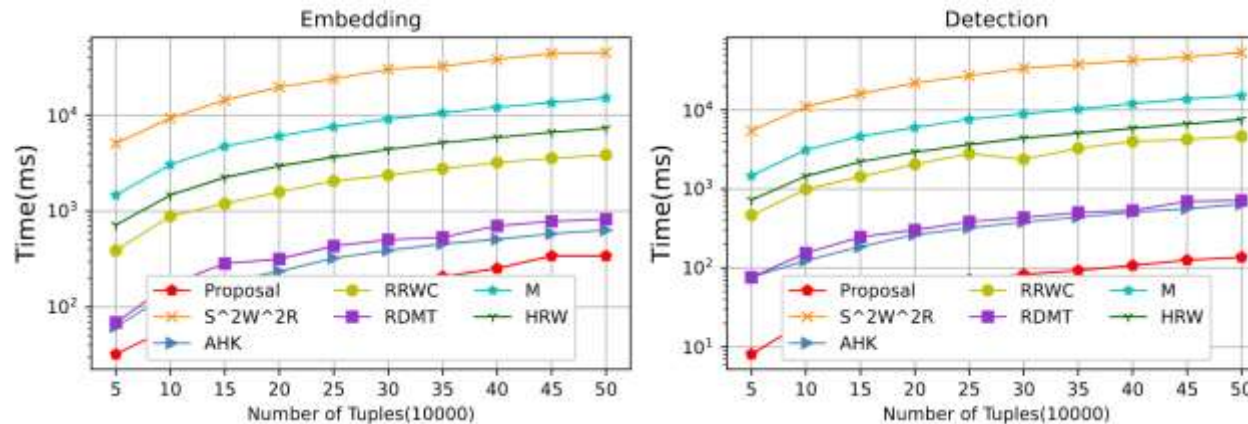
Experiments

Usability Verification

Data Distortions

Attribute	$\Delta\mu$	$\Delta\sigma$
Elevation	-1.89e-4	2.33e-4
Slope	-1.6e-4	7.98e-05
Horizontal_Distance_To_Hydrology	1.31e-4	-2.57e-4
Vertical_Distance_To_Hydrology	1.39e-4	-6.24e-6
Horizontal_Distance_To_Roadways	1.08e-4	-4.66e-7
Hillshade_9am	8.9e-5	8.56e-5
Hillshade_3pm	-9.8e-5	2.78e-5
Horizontal_Distance_To_Fire_Points	4.6e-5	1.81e-4

Overhead



Conclusion

- **Introduce Attribute column attacks and mix-match attribute attacks.**
- **Train an attribute classifier for resisting these attacks.**
- **Design a high-synchronization and robust virtual primary key scheme for watermarking relational databases.**
- **Performance experiments substantiate our scheme.**

Thanks

Any concern feel free to reach me out:
yangke@iie.ac.cn